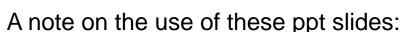
Chapter 6 Wireless and Mobile Networks

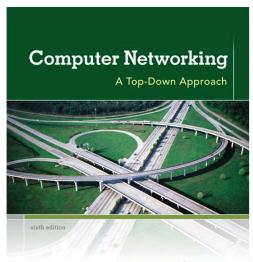


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KUROSE ROSS

Computer Networking: A Top Down Approach 6th edition Jim Kurose, Keith Ross Addison-Wesley March 2012

Ch. 6: Wireless and Mobile Networks

Background:

- # wireless (mobile) phone subscribers now exceeds # wired phone subscribers (5-to-I)!
- # wireless Internet-connected devices equals # wireline Internet-connected devices
 - laptops, Internet-enabled phones promise anytime untethered Internet access
- two important (but different) challenges
 - wireless: communication over wireless link
 - mobility: handling the mobile user who changes point of attachment to network

Chapter 6 outline

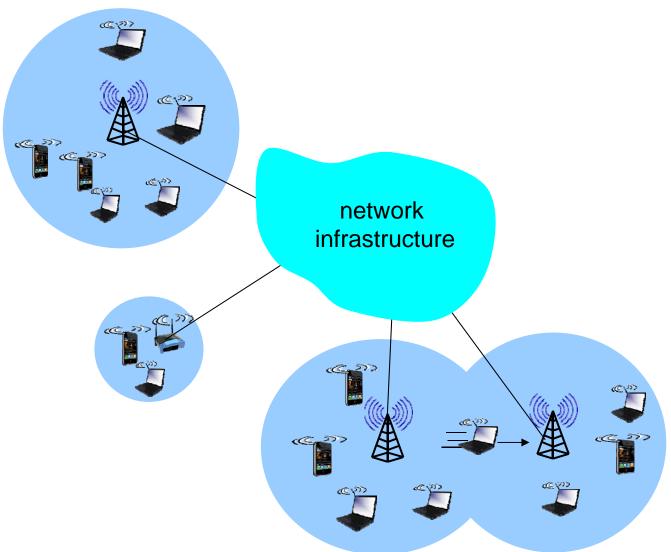
6.1 Introduction

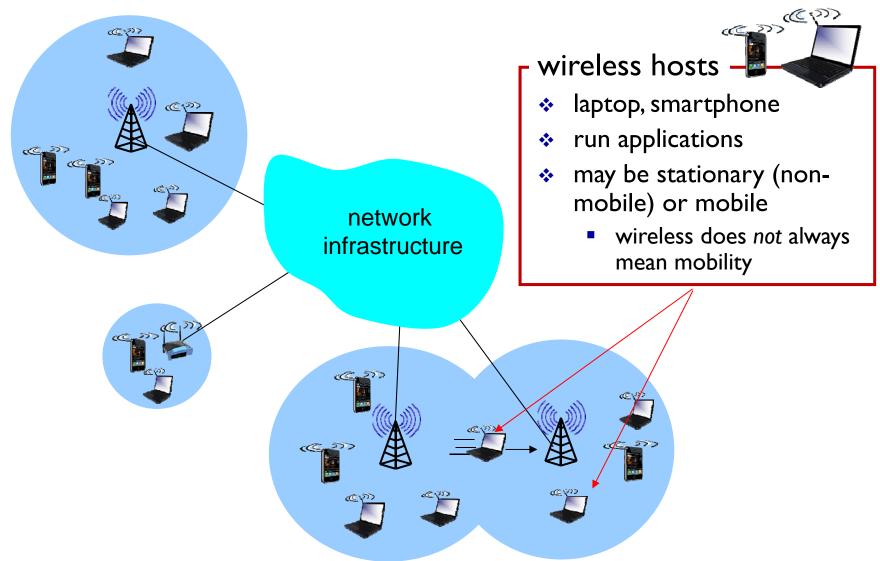
Wireless

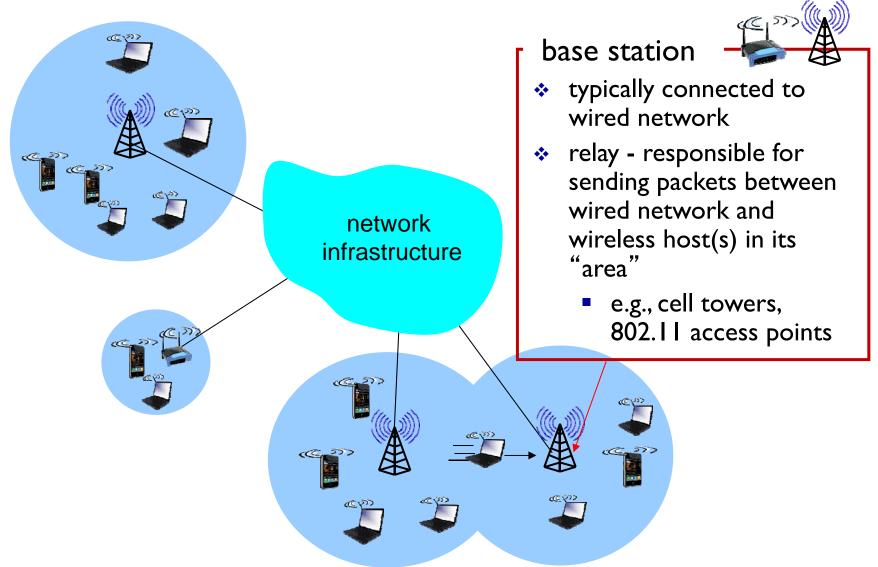
- 6.2 Wireless links, characteristics
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- 6.3 IEEE 802.11 wireless LANs ("Wi-Fi")
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 - standards (e.g., GSM)

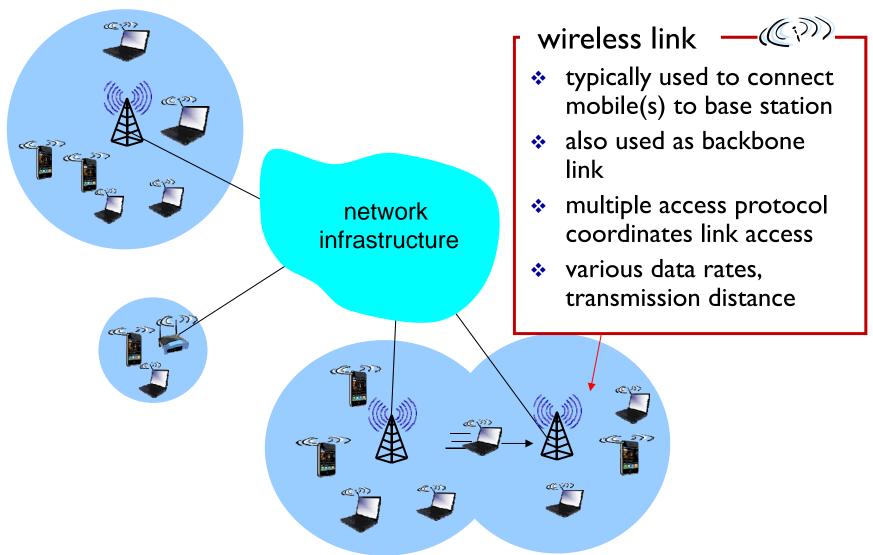
Mobility

- 6.5 Principles: addressing and routing to mobile users
- 6.6 Mobile IP
- 6.7 Handling mobility in cellular networks
- 6.8 Mobility and higher-layer protocols
- 6.9 Summary

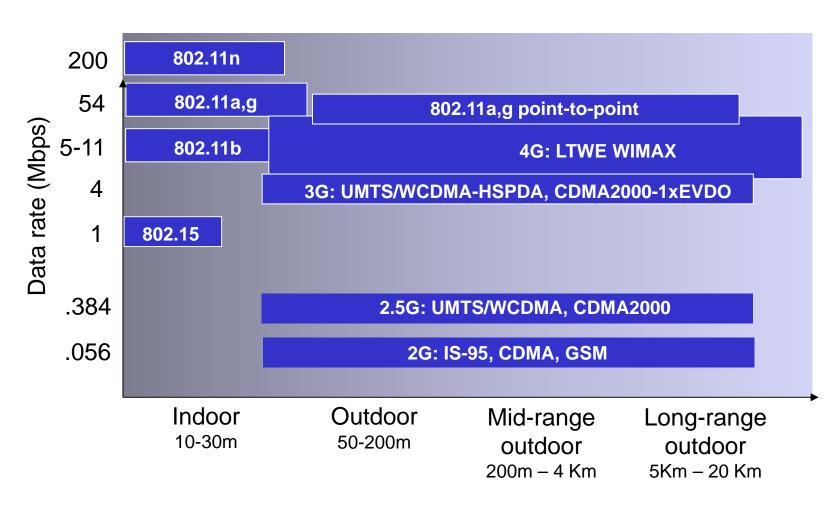


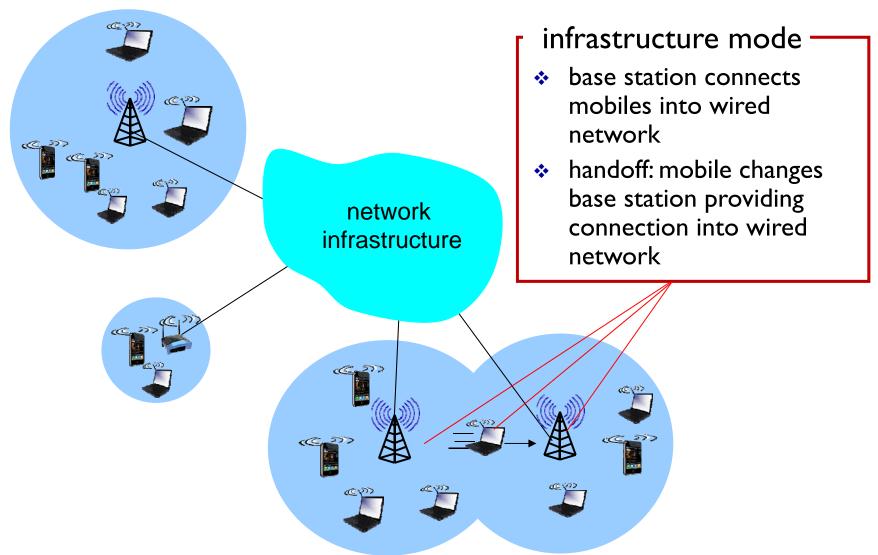


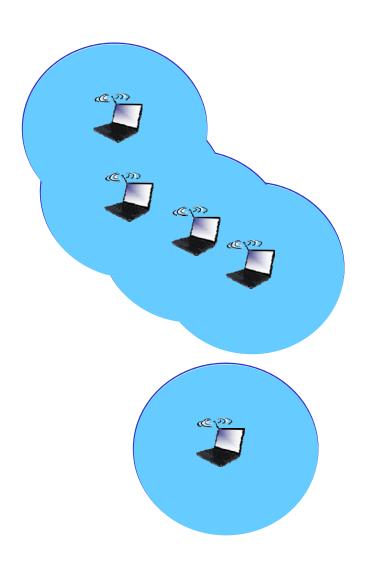




Characteristics of selected wireless links







ad hoc mode

- no base stations
- nodes can only transmit to other nodes within link coverage
- nodes organize themselves into a network: route among themselves

Wireless network taxonomy

	single hop	multiple hops
infrastructure (e.g., APs)	host connects to base station (WiFi, WiMAX, cellular) which connects to larger Internet	host may have to relay through several wireless nodes to connect to larger Internet: mesh net
no infrastructure	no base station, no connection to larger Internet (Bluetooth, ad hoc nets)	no base station, no connection to larger Internet. May have to relay to reach other a given wireless node MANET, VANET

Chapter 6 outline

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Wireless Link Characteristics (I)

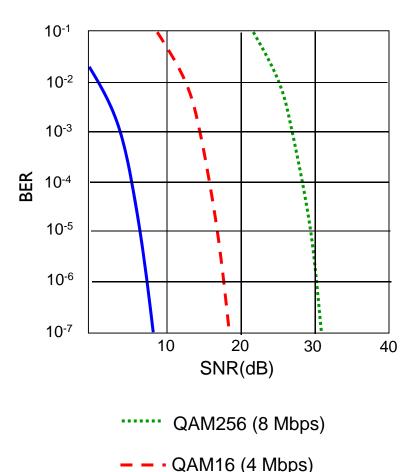
important differences from wired link

- decreased signal strength: radio signal attenuates as it propagates through matter (path loss)
- interference from other sources: standardized wireless network frequencies (e.g., 2.4 GHz) shared by other devices (e.g., phone); devices (motors) interfere as well
- multipath propagation: radio signal reflects off objects ground, arriving ad destination at slightly different times

.... make communication across (even a point to point) wireless link much more "difficult"

Wireless Link Characteristics (2)

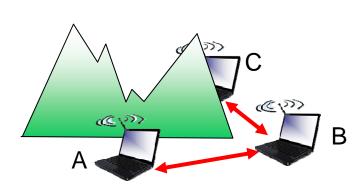
- SNR: signal-to-noise ratio
 - larger SNR easier to extract signal from noise (a "good thing")
- SNR versus BER tradeoffs
 - given physical layer: increase
 power -> increase SNR >decrease BER
 - given SNR: choose physical layer that meets BER requirement, giving highest thruput
 - SNR may change with mobility: dynamically adapt physical layer (modulation technique, rate)



BPSK (1 Mbps)

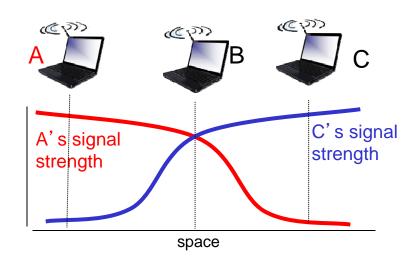
Wireless network characteristics

Multiple wireless senders and receivers create additional problems (beyond multiple access):



Hidden terminal problem

- ♣ B,A hear each other
- B, C hear each other
- A, C can not hear each other means A, C unaware of their interference at B



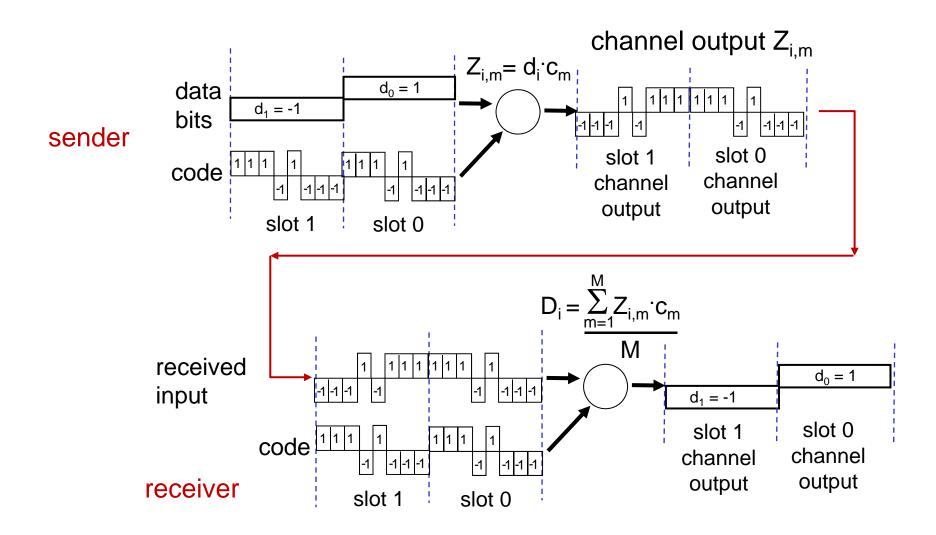
Signal attenuation:

- B,A hear each other
- * B, C hear each other
- A, C can not hear each other interfering at B

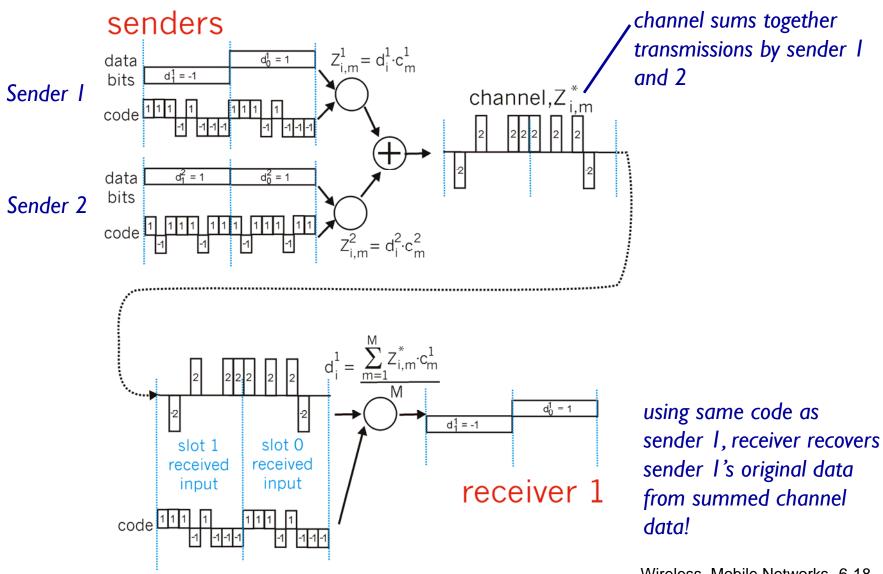
Code Division Multiple Access (CDMA)

- unique "code" assigned to each user; i.e., code set partitioning
 - all users share same frequency, but each user has own "chipping" sequence (i.e., code) to encode data
 - allows multiple users to "coexist" and transmit simultaneously with minimal interference (if codes are "orthogonal")
- encoded signal = (original data) X (chipping sequence)
- decoding: inner-product of encoded signal and chipping sequence

CDMA encode/decode



CDMA: two-sender interference



Wireless, Mobile Networks 6-18

Chapter 6 outline

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IEEE 802.11 Wireless LAN

802.11b

- 2.4-5 GHz unlicensed spectrum
- up to 11 Mbps
- direct sequence spread spectrum (DSSS) in physical layer
 - all hosts use same chipping code

802.11a

- 5-6 GHz range
- up to 54 Mbps

802.11g

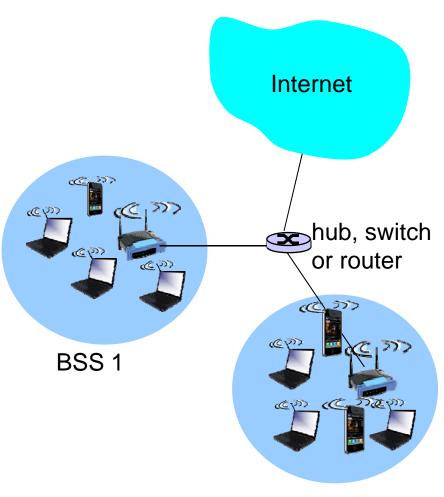
- 2.4-5 GHz range
- up to 54 Mbps

802. I In: multiple antennae

- 2.4-5 GHz range
- up to 200 Mbps

- all use CSMA/CA for multiple access
- all have base-station and ad-hoc network versions

802.11 LAN architecture



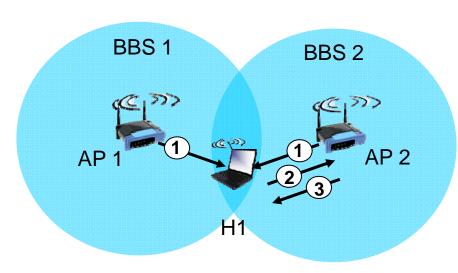
BSS 2

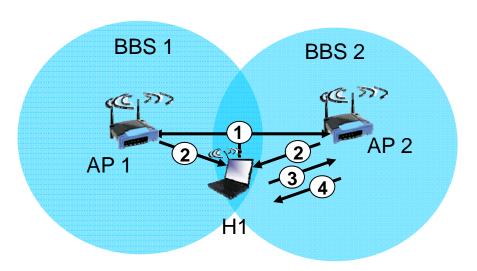
- wireless host communicates with base station
 - base station = access point (AP)
- * Basic Service Set (BSS) (aka "cell") in infrastructure mode contains:
 - wireless hosts
 - access point (AP): base station
 - ad hoc mode: hosts only

802. I I: Channels, association

- ❖ 802.11b: 2.4GHz-2.485GHz spectrum divided into 11 channels at different frequencies
 - AP admin chooses frequency for AP
 - interference possible: channel can be same as that chosen by neighboring AP!
- host: must associate with an AP
 - scans channels, listening for beacon frames containing AP's name (SSID) and MAC address
 - selects AP to associate with
 - may perform authentication [Chapter 8]
 - will typically run DHCP to get IP address in AP's subnet

802.11: passive/active scanning





passive scanning:

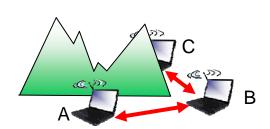
- (I) beacon frames sent from APs
- (2) association Request frame sent: H1 to selected AP
- (3) association Response frame sent from selected AP to HI

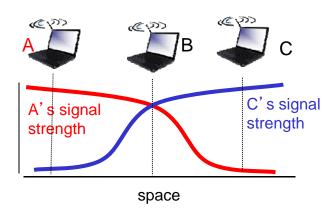
active scanning:

- (1) Probe Request frame broadcast from H1
- (2) Probe Response frames sent from APs
- (3) Association Request frame sent: H1 to selected AP
- (4) Association Response frame sent from selected AP to H1

IEEE 802.11: multiple access

- avoid collisions: 2⁺ nodes transmitting at same time
- ❖ 802.11: CSMA sense before transmitting
 - don't collide with ongoing transmission by other node
- * 802.11: no collision detection!
 - difficult to receive (sense collisions) when transmitting due to weak received signals (fading)
 - can't sense all collisions in any case: hidden terminal, fading
 - goal: avoid collisions: CSMA/C(ollision)A(voidance)





IEEE 802.11 MAC Protocol: CSMA/CA

802.11 sender

1 if sense channel idle for **DIFS** then transmit entire frame (no CD)

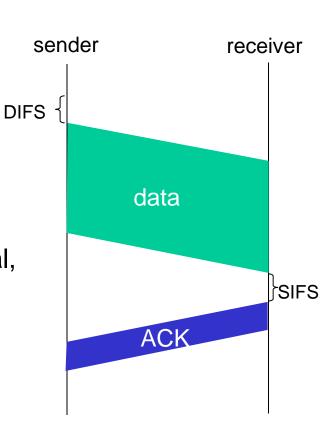
2 if sense channel busy then

start random backoff time
timer counts down while channel idle
transmit when timer expires
if no ACK, increase random backoff interval,
repeat 2

<u>802.11 receiver</u>

- if frame received OK

return ACK after **SIFS** (ACK needed due to hidden terminal problem)

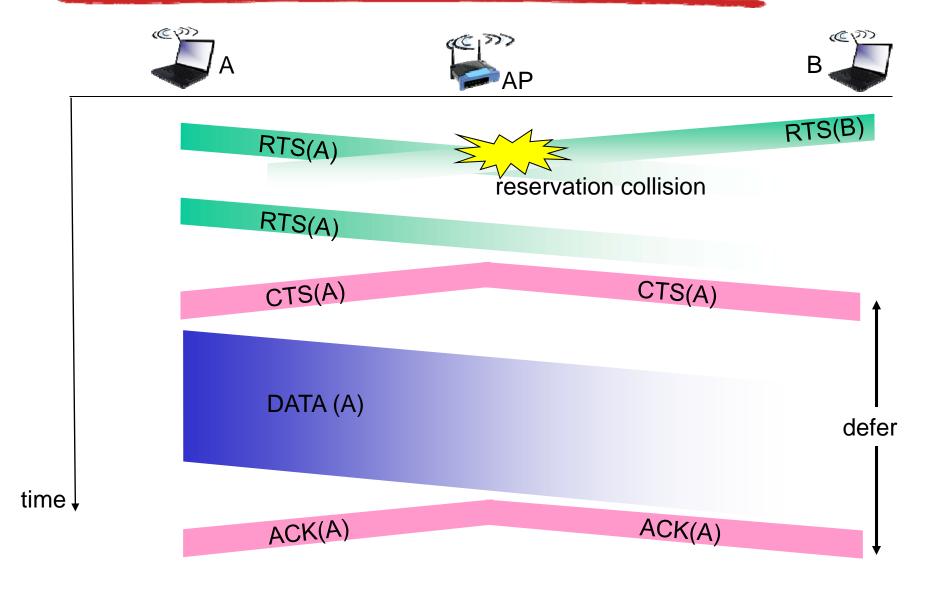


Avoiding collisions (more)

- idea: allow sender to "reserve" channel rather than random access of data frames: avoid collisions of long data frames
- sender first transmits small request-to-send (RTS) packets to BS using CSMA
 - RTSs may still collide with each other (but they' re short)
- BS broadcasts clear-to-send CTS in response to RTS
- CTS heard by all nodes
 - sender transmits data frame
 - other stations defer transmissions

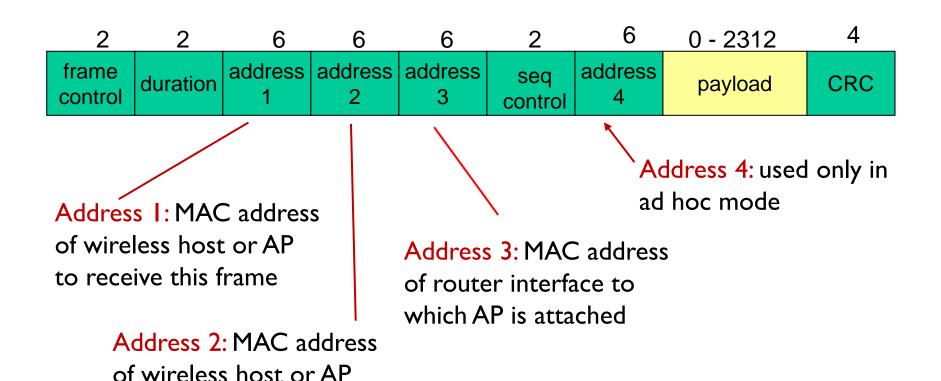
avoid data frame collisions completely using small reservation packets!

Collision Avoidance: RTS-CTS exchange

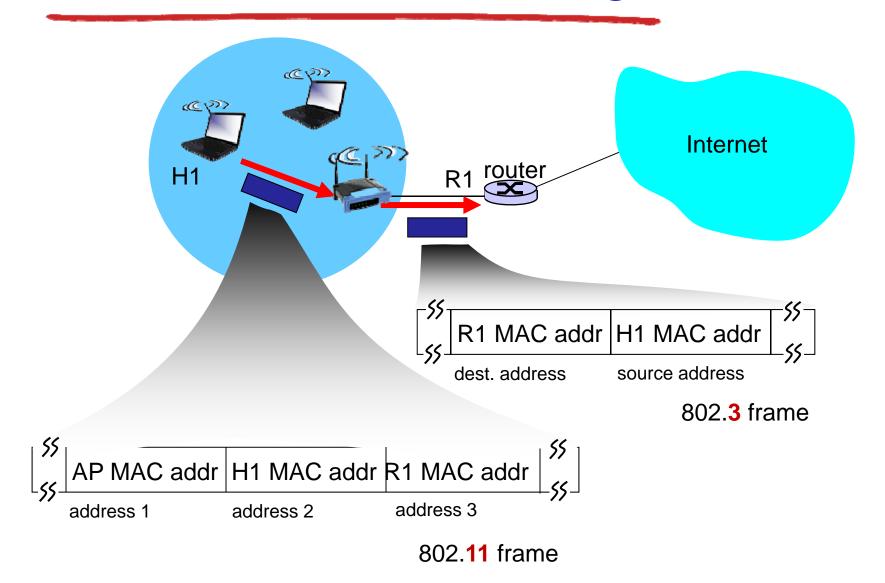


802.11 frame: addressing

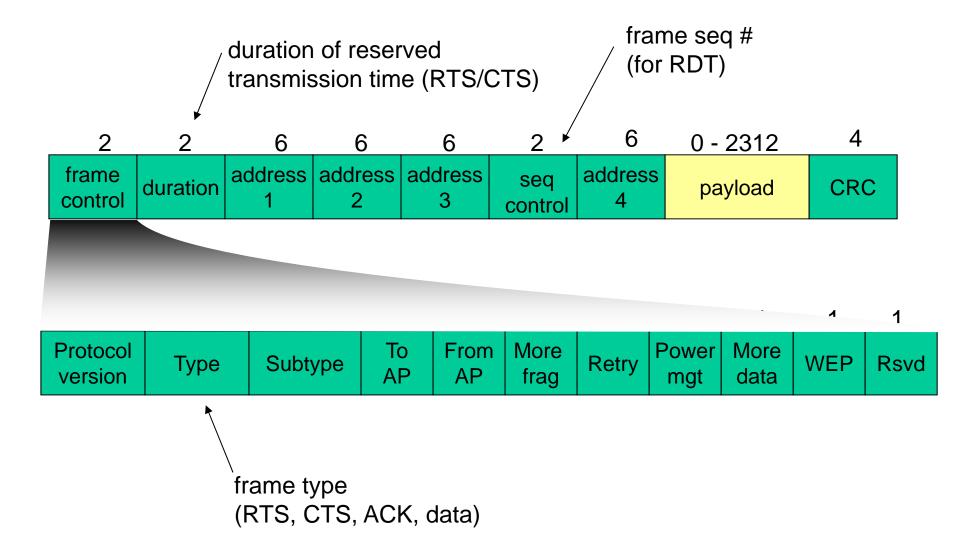
transmitting this frame



802.11 frame: addressing

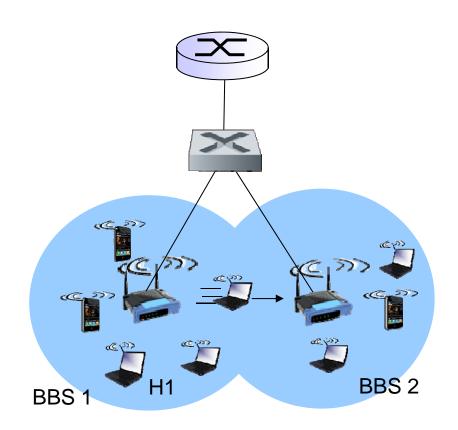


802.11 frame: more



802.11: mobility within same subnet

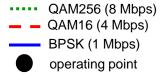
- HI remains in same IP subnet: IP address can remain same
- switch: which AP is associated with HI?
 - self-learning (Ch. 5): switch will see frame from HI and "remember" which switch port can be used to reach HI

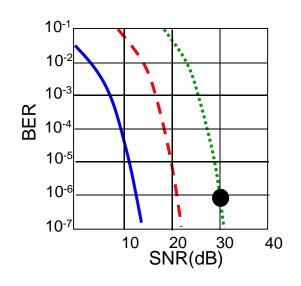


802. I I: advanced capabilities

Rate adaptation

base station, mobile
 dynamically change
 transmission rate
 (physical layer modulation
 technique) as mobile
 moves, SNR varies





- 1. SNR decreases, BER increase as node moves away from base station
- 2. When BER becomes too high, switch to lower transmission rate but with lower BER

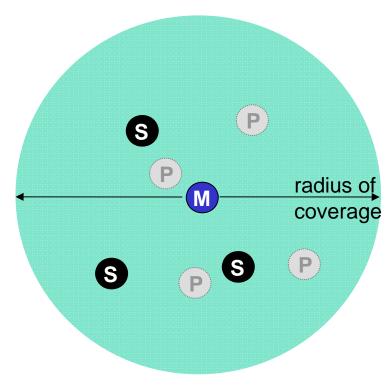
802. I I: advanced capabilities

power management

- node-to-AP: "I am going to sleep until next beacon frame"
 - AP knows not to transmit frames to this node
 - node wakes up before next beacon frame
- beacon frame: contains list of mobiles with APto-mobile frames waiting to be sent
 - node will stay awake if AP-to-mobile frames to be sent; otherwise sleep again until next beacon frame

802.15: personal area network

- less than 10 m diameter
- replacement for cables (mouse, keyboard, headphones)
- ad hoc: no infrastructure
- master/slaves:
 - slaves request permission to send (to master)
 - master grants requests
- * 802.15: evolved from Bluetooth specification
 - 2.4-2.5 GHz radio band
 - up to 721 kbps



- Master device
- S Slave device
- P Parked device (inactive)

Chapter 6 outline

6. I Introduction

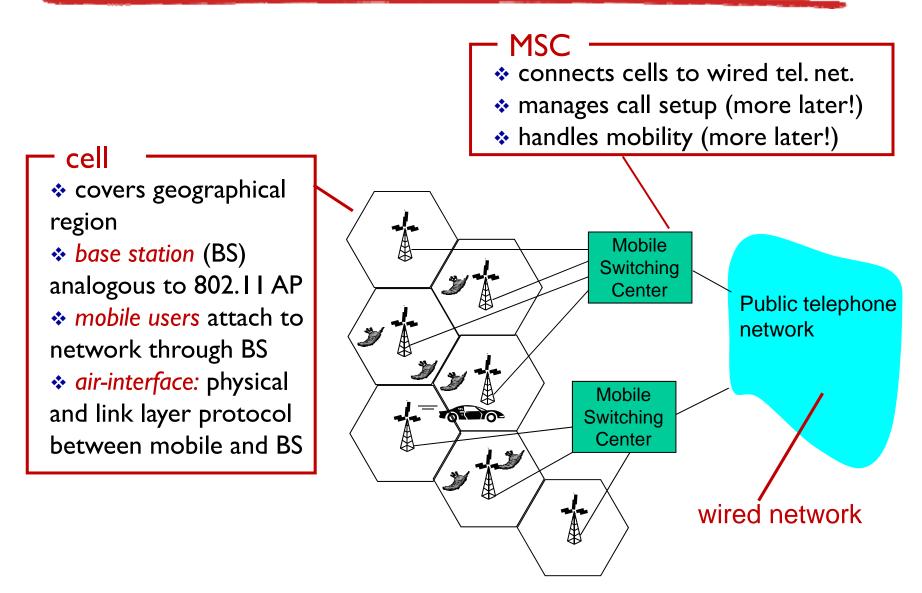
Wireless

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 - standards (e.g., GSM)

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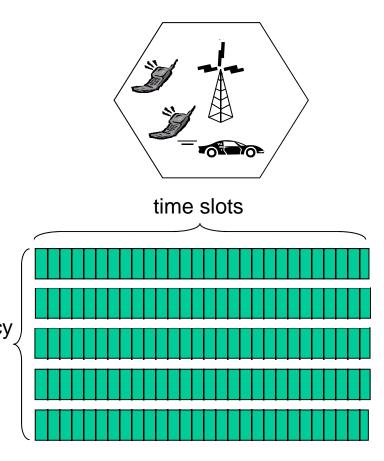
Components of cellular network architecture



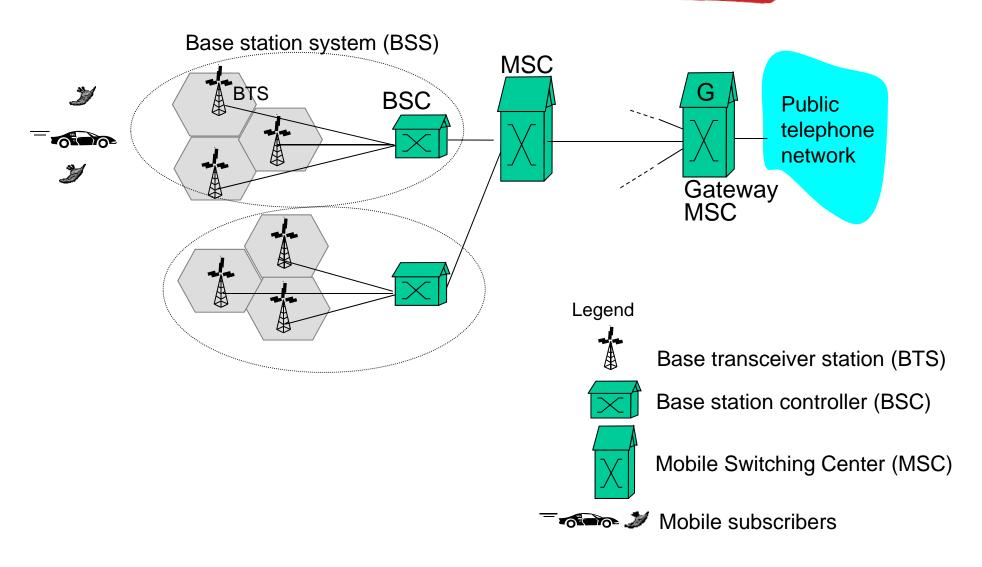
Cellular networks: the first hop

Two techniques for sharing mobile-to-BS radio spectrum

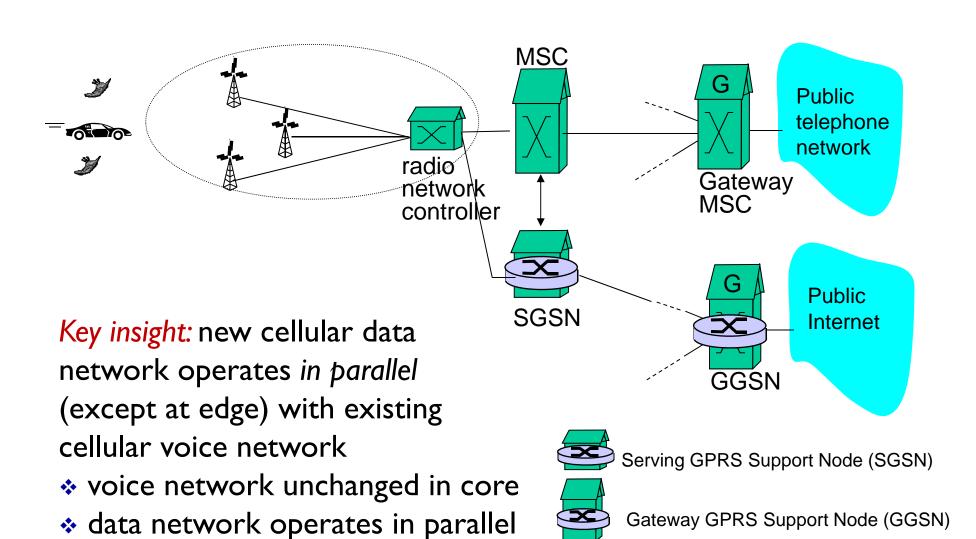
- combined FDMA/TDMA: divide spectrum in frequency channels, divide each channel into time slots
- CDMA: code division multiple access
 frequency bands



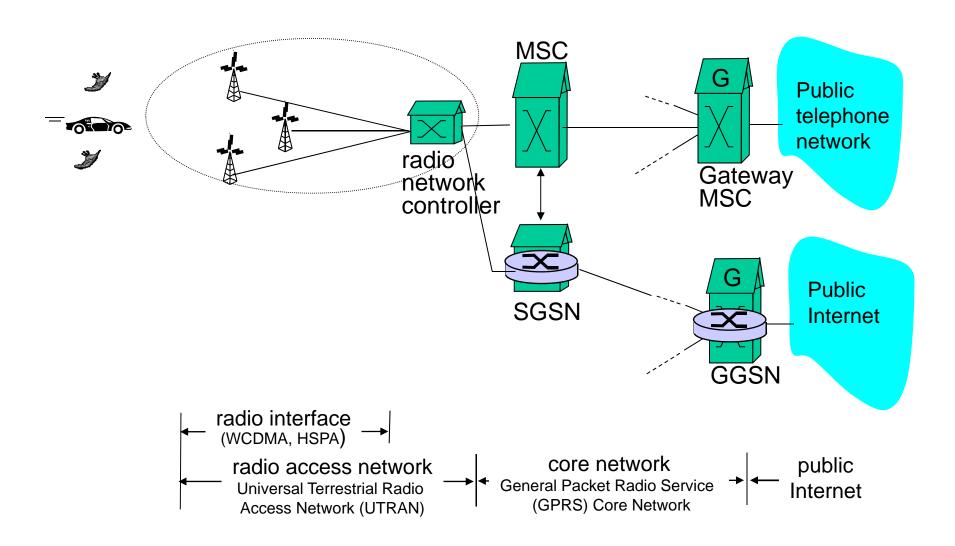
2G (voice) network architecture



3G (voice+data) network architecture



3G (voice+data) network architecture



Chapter 6 outline

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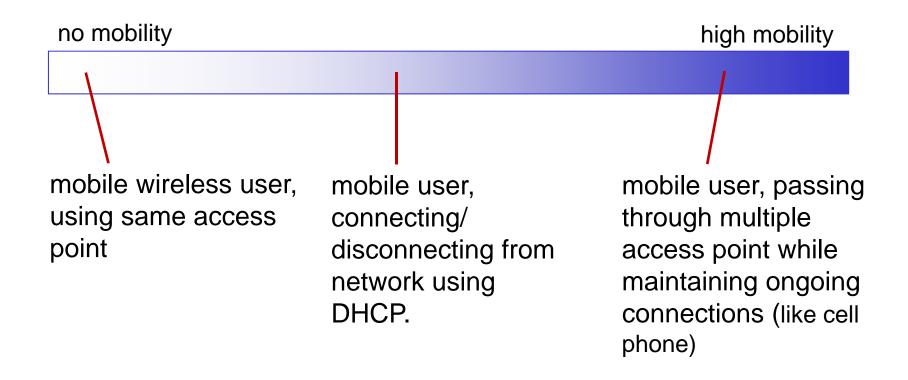
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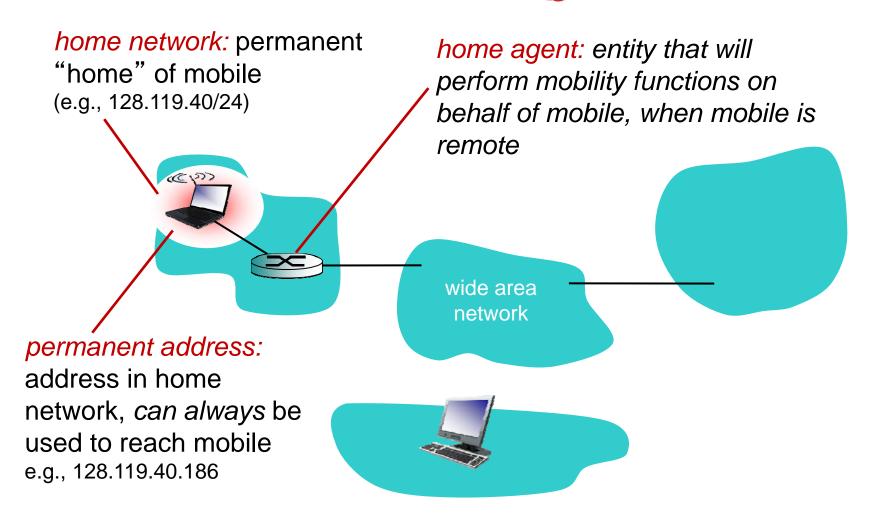
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What is mobility?

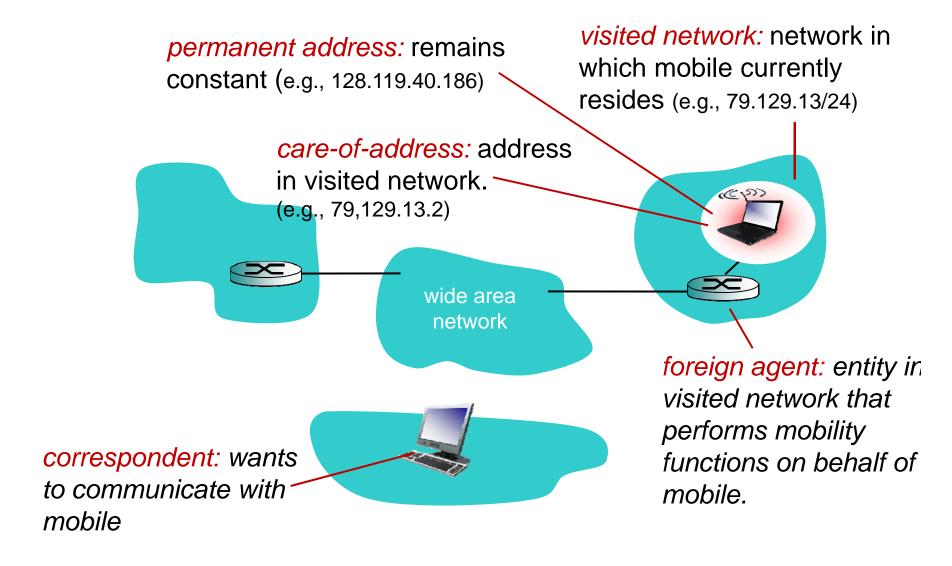
spectrum of mobility, from the network perspective:



Mobility: vocabulary



Mobility: more vocabulary



How do you contact a mobile friend:

Consider friend frequently changing addresses, how do you find her?

search all phone books?

call her parents?

expect her to let you know where he/she is?

I wonder where Alice moved to?



Mobility: approaches

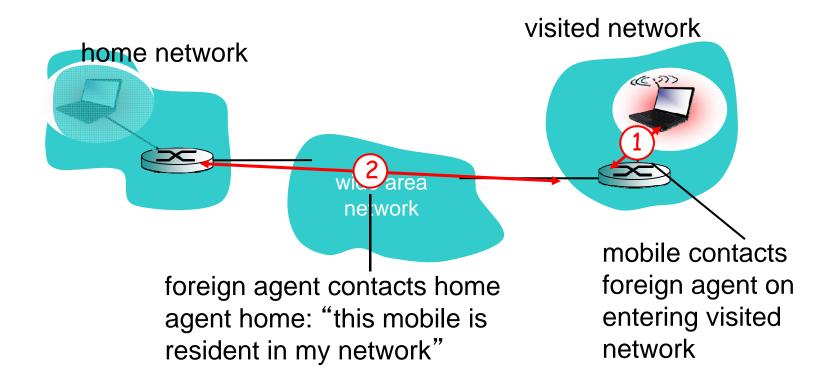
- let routing handle it: routers advertise permanent address of mobile-nodes-in-residence via usual routing table exchange.
 - routing tables indicate where each mobile located
 - no changes to end-systems
- let end-systems handle it:
 - indirect routing: communication from correspondent to mobile goes through home agent, then forwarded to remote
 - direct routing: correspondent gets foreign address of mobile, sends directly to mobile

Mobility: approaches

- let routing handle it: r vertise permanent address of mobile-nodes-in-r sual routing table exchange. not

 - no changes to
 - routing tables scalable to millions of each mobile located
 - mobiles
- let end-systems handle it.
 - indirect routing: communication from correspondent to mobile goes through home agent, then forwarded to remote
 - direct routing: correspondent gets foreign address of mobile, sends directly to mobile

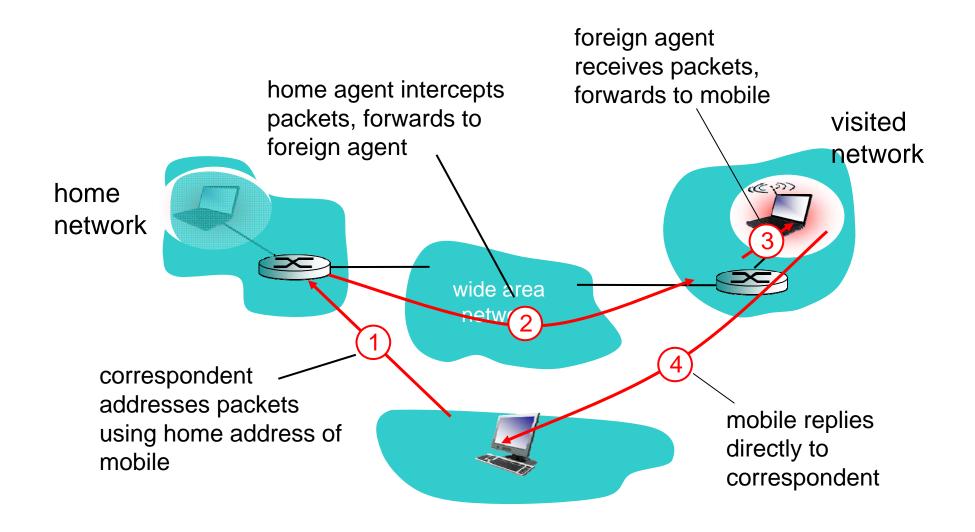
Mobility: registration



end result:

- foreign agent knows about mobile
- home agent knows location of mobile

Mobility via indirect routing



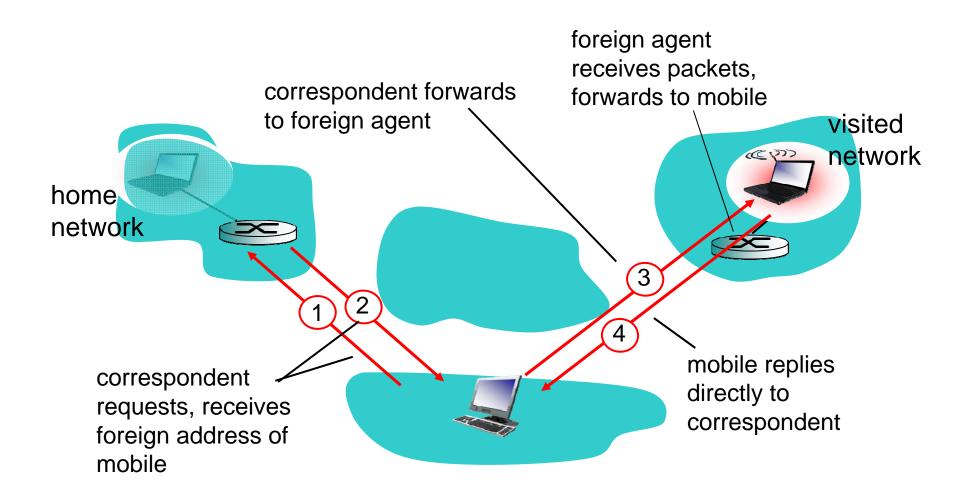
Indirect Routing: comments

- mobile uses two addresses:
 - permanent address: used by correspondent (hence mobile location is transparent to correspondent)
 - care-of-address: used by home agent to forward datagrams to mobile
- foreign agent functions may be done by mobile itself
- triangle routing: correspondent-home-network-mobile
 - inefficient when correspondent, mobile are in same network

Indirect routing: moving between networks

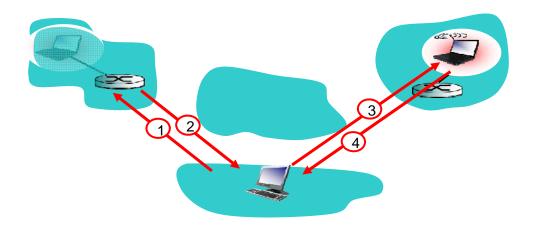
- suppose mobile user moves to another network
 - registers with new foreign agent
 - new foreign agent registers with home agent
 - home agent update care-of-address for mobile
 - packets continue to be forwarded to mobile (but with new care-of-address)
- mobility, changing foreign networks transparent: on going connections can be maintained!

Mobility via direct routing



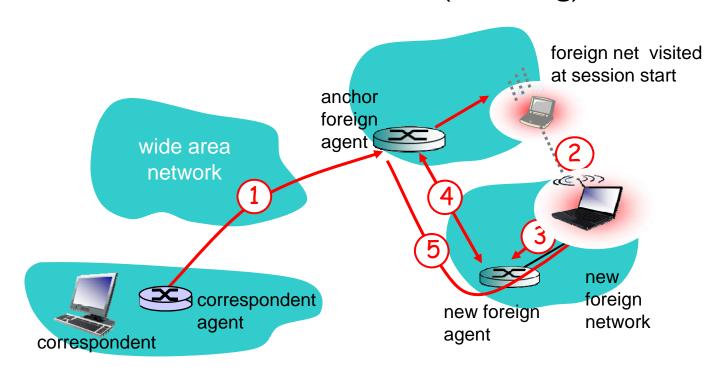
Mobility via direct routing: comments

- overcome triangle routing problem
- non-transparent to correspondent: correspondent must get care-of-address from home agent
 - what if mobile changes visited network?



Accommodating mobility with direct routing

- anchor foreign agent: FA in first visited network
- data always routed first to anchor FA
- when mobile moves: new FA arranges to have data forwarded from old FA (chaining)



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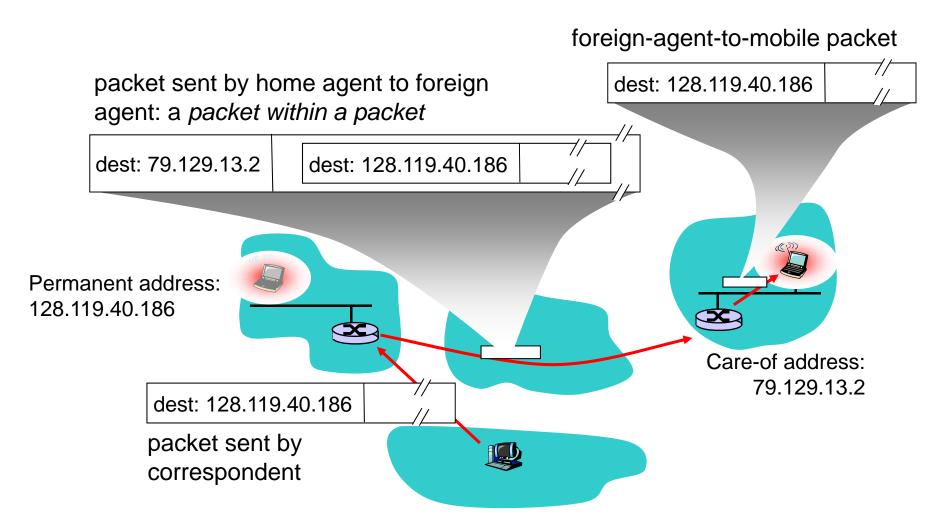
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Mobile IP

- * RFC 3344
- * has many features we've seen:
 - home agents, foreign agents, foreign-agent registration, care-of-addresses, encapsulation (packet-within-apacket)
- three components to standard:
 - indirect routing of datagrams
 - agent discovery
 - registration with home agent

Mobile IP: indirect routing

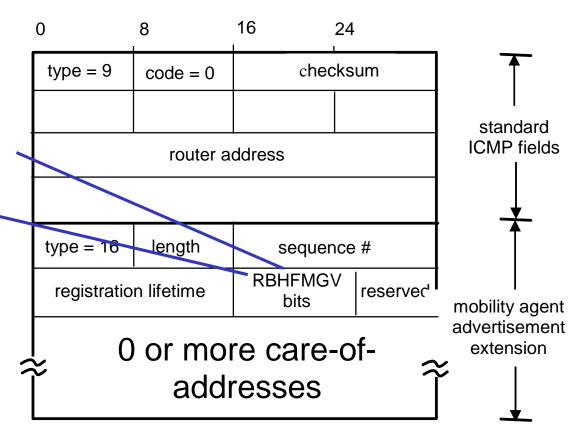


Mobile IP: agent discovery

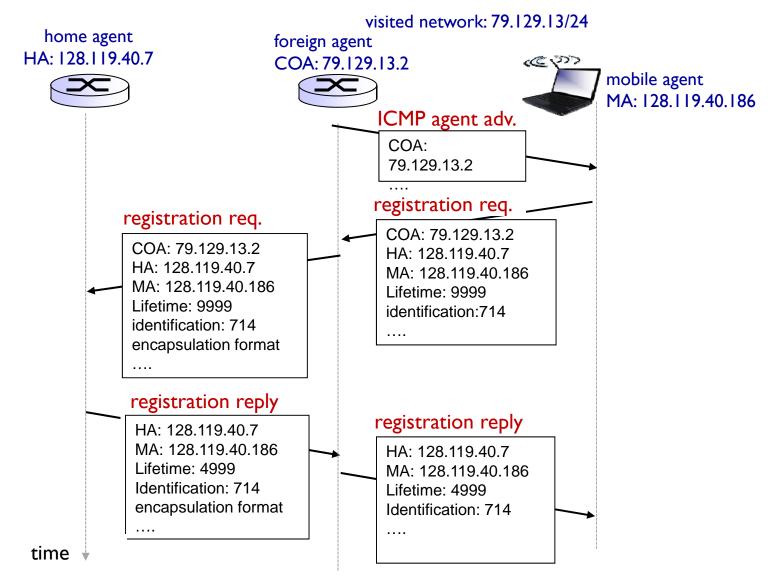
agent advertisement: foreign/home agents advertise service by broadcasting ICMP messages (typefield = 9)

H,F bits: home and/or foreign agent

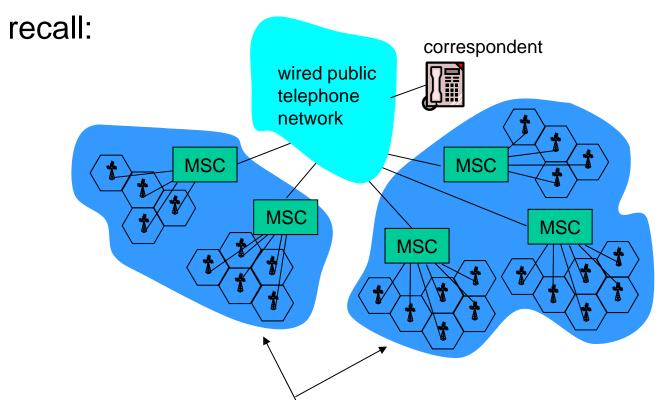
R bit: registration required



Mobile IP: registration example



Components of cellular network architecture

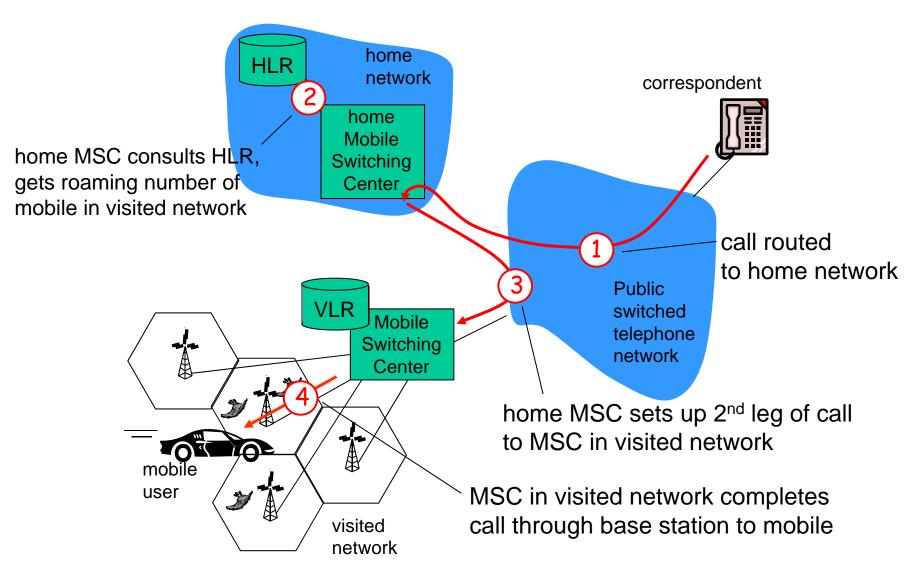


different cellular networks, operated by different providers

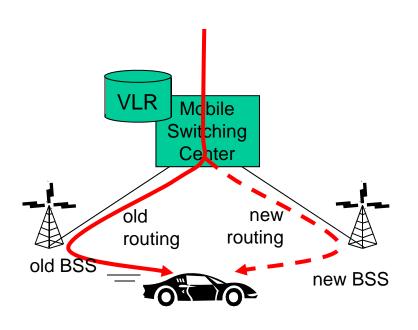
Handling mobility in cellular networks

- home network: network of cellular provider you subscribe to (e.g., Sprint PCS, Verizon)
 - home location register (HLR): database in home network containing permanent cell phone #, profile information (services, preferences, billing), information about current location (could be in another network)
- visited network: network in which mobile currently resides
 - visitor location register (VLR): database with entry for each user currently in network
 - could be home network

GSM: indirect routing to mobile

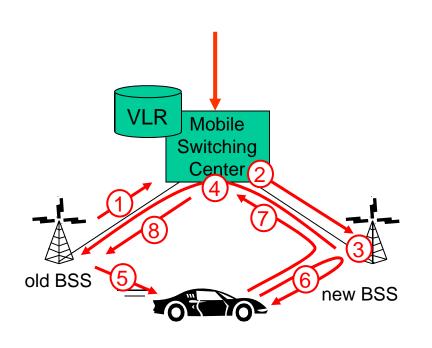


GSM: handoff with common MSC



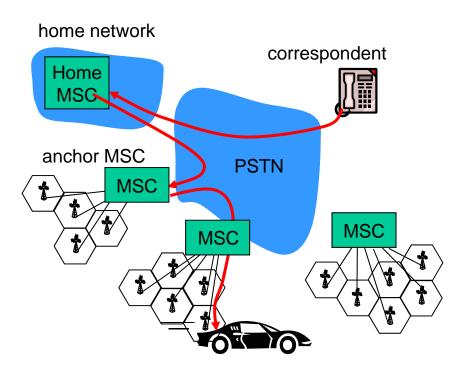
- handoff goal: route call via new base station (without interruption)
- reasons for handoff:
 - stronger signal to/from new BSS (continuing connectivity, less battery drain)
 - load balance: free up channel in current BSS
 - GSM doesnt mandate why to perform handoff (policy), only how (mechanism)
- handoff initiated by old BSS

GSM: handoff with common MSC



- 1. old BSS informs MSC of impending handoff, provides list of 1+ new BSSs
- 2. MSC sets up path (allocates resources) to new BSS
- 3. new BSS allocates radio channel for use by mobile
- 4. new BSS signals MSC, old BSS: ready
- 5. old BSS tells mobile: perform handoff to new BSS
- 6. mobile, new BSS signal to activate new channel
- 7. mobile signals via new BSS to MSC: handoff complete. MSC reroutes call
- 8 MSC-old-BSS resources released

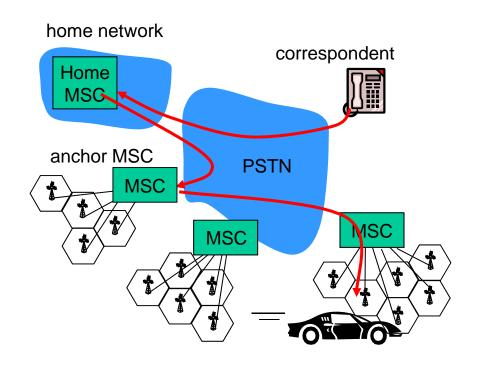
GSM: handoff between MSCs



(a) before handoff

- anchor MSC: first MSC visited during call
 - call remains routed through anchor MSC
- new MSCs add on to end of MSC chain as mobile moves to new MSC
- optional path minimization step to shorten multi-MSC chain

GSM: handoff between MSCs



(b) after handoff

- anchor MSC: first MSC visited during call
 - call remains routed through anchor MSC
- new MSCs add on to end of MSC chain as mobile moves to new MSC
- optional path minimization step to shorten multi-MSC chain

Mobility: GSM versus Mobile IP

GSM element	Comment on GSM element Mo	obile IP element
Home system	Network to which mobile user's permanent phone number belongs	Home network
Gateway Mobile Switching Center, or "home MSC". Home Location Register (HLR)	Home MSC: point of contact to obtain routable address of mobile user. HLR: database in home system containing permanent phone number, profile information, current location of mobile user, subscription information	Home agent
Visited System	Network other than home system where mobile user is currently residing	Visited network
Visited Mobile services Switching Center. Visitor Location Record (VLR)	Visited MSC: responsible for setting up calls to/from mobile nodes in cells associated with MSC. VLR: temporary database entry in visited system, containing subscription information for each visiting mobile user	Foreign agent
Mobile Station Roaming Number (MSRN), or "roaming number"	Routable address for telephone call segment between home MSC and visited MSC, visible to neither the mobile nor the correspondent.	Care-of- address

Wireless, mobility: impact on higher layer protocols

- logically, impact should be minimal ...
 - best effort service model remains unchanged
 - TCP and UDP can (and do) run over wireless, mobile
- ... but performance-wise:
 - packet loss/delay due to bit-errors (discarded packets, delays for link-layer retransmissions), and handoff
 - TCP interprets loss as congestion, will decrease congestion window un-necessarily
 - delay impairments for real-time traffic
 - limited bandwidth of wireless links

Chapter 6 summary

Wireless

- wireless links:
 - capacity, distance
 - channel impairments
 - CDMA
- ❖ IEEE 802.11 ("Wi-Fi")
 - CSMA/CA reflects wireless channel characteristics
- cellular access
 - architecture
 - standards (e.g., GSM, 3G, 4G LTE)

Mobility

- principles: addressing, routing to mobile users
 - home, visited networks
 - direct, indirect routing
 - care-of-addresses
- case studies
 - mobile IP
 - mobility in GSM
- impact on higher-layer protocols